

MINIMAL NORMAL AND COMMUTING COMPLETIONS

David P. Kimsey and Hugo J. Woerdeman *

Department of Mathematics

Drexel University

Philadelphia, PA 19104

Abstract

We study the minimal normal completion problem: given $A \in \mathbb{C}^{n \times n}$, how do we find a $(n + q) \times (n + q)$ normal matrix $A_{ext} := \begin{pmatrix} A & A_{12} \\ A_{21} & A_{22} \end{pmatrix}$ of smallest possible size? We will show that this smallest number q of rows and columns we need to add, called the *normal defect* of A , satisfies

$$\text{nd}(A) \geq \max\{i_-(AA^* - A^*A), i_+(AA^* - A^*A)\},$$

where $i_{\pm}(M)$ denotes the number of positive and negative eigenvalues of the Hermitian matrix M . Subsequently, we will show that for some matrices a minimal normal completion can be chosen to be a multiple of a unitary, addressing a conjecture from [H. J. Woerdeman, *Linear and Multilinear Algebra* 36 (1993), 59–68].

In addition, we study the related question where $A \in \mathbb{C}^{n \times n}$ and $B \in \mathbb{C}^{n \times n}$ are given, and where we look for $A_{ext} := \begin{pmatrix} A & A_{12} \\ A_{21} & A_{22} \end{pmatrix}$ and $B_{ext} := \begin{pmatrix} B & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$ such they commute and are of smallest possible size.

*This work was partially supported by NSF grant DMS-0500678. David P. Kimsey performed the research as part of an REU project.

1 INTRODUCTION

The Minimal Normal Completion problem was introduced in [7] and concerns the following. Given $A \in \mathbb{C}^{n \times n}$, we wish to find a smallest possible normal matrix with A as a principal submatrix. Recall that a matrix A is normal if and only if the *commutator* of A and its conjugate transpose A^* , denoted by $[A, A^*] := AA^* - A^*A$, equals 0. In other words, we would like to find a normal completion of

$$\begin{pmatrix} A & ? \\ ? & ? \end{pmatrix} : \begin{matrix} \mathbb{C}^n \\ \oplus \\ \mathbb{C}^q \end{matrix} \rightarrow \begin{matrix} \mathbb{C}^n \\ \oplus \\ \mathbb{C}^q \end{matrix}$$

of smallest possible size (thus smallest possible q). We shall call this smallest number q the *normal defect* of A , and denote it by $\text{nd}(A)$. Clearly, $\text{nd}(A) = 0$ if and only if A is normal. As observed in [2], the matrix $\begin{pmatrix} A & A^* \\ A^* & A \end{pmatrix}$ is normal, so it follows that for a $n \times n$ matrix A we have that $\text{nd}(A) \leq n$. As was observed in [7], and as we shall see further on, we have in fact that $\text{nd}(A) \leq n - 1$. It is also not hard to come up with the lower bound $\text{nd}(A) \geq \frac{1}{2}\text{rank}(AA^* - A^*A)$. Indeed if $\begin{pmatrix} A & A_{12} \\ A_{21} & A_{22} \end{pmatrix}$ is of size $(n + q) \times (n + q)$ and normal then $AA^* - A^*A = A_{21}^*A_{21} - A_{12}A_{12}^*$ and thus $\text{rank}(AA^* - A^*A) \leq \text{rank}(A_{21}^*A_{21}) + \text{rank}(A_{12}A_{12}^*) \leq q + q = 2q$.

In order to obtain sharper bounds for $\text{nd}(A)$, the so-called *unitary defect* was introduced in [7]. It corresponds to the smallest number of rows and columns we need to add to A such that the completion is a multiple of a unitary matrix, and as it turns out we have that

$$\text{ud}(A) := \text{rank}(\|A\|^2 - A^*A), \quad (1.1)$$

where $\|\cdot\|$ denotes the spectral norm. As multiples of unitaries are normal we clearly have that $\text{nd}(A) \leq \text{ud}(A)$. Formula (1.1) implies that $\text{ud}(A) \leq n - 1$, which yields $\text{nd}(A) \leq n - 1$.

In order to state a conjecture from [7], let us recall that a matrix $A \in \mathbb{C}^{n \times n}$ is called *unitarily*

*Keywords and phrases: commuting completions, normal completions, normal defect, inertia, inverse defect, unitary defect.

*2000 AMS subject classification: 15A57, 15A24, 15A42.

reducible if $A = U^* \begin{pmatrix} A_1 & 0 \\ 0 & A_2 \end{pmatrix} U$, with U unitary and A_1, A_2 of nontrivial size. Clearly, with A as above we have that

$$\text{nd}(A) \leq \text{nd}(A_1) + \text{nd}(A_2) \leq \text{ud}(A_1) + \text{ud}(A_2) \leq \text{ud}(A). \quad (1.2)$$

As soon as $\|A_1\| \neq \|A_2\|$ we have that the last inequality in (1.2) is strict, and thus $\text{nd}(A) < \text{ud}(A)$ in that case. So for a general statement for the case when $\text{nd}(A) = \text{ud}(A)$, it is natural to require that A is *unitarily irreducible*, which by definition means that A is not unitarily reducible. An open question from [7] is whether the following conjecture holds.

Conjecture 1.1 *For an unitarily irreducible matrix A we have that $\text{nd}(A) = \text{ud}(A)$.*

In this paper we refine some of the estimates for $\text{nd}(A)$ and as a result obtain more evidence for this conjecture. Let us mention that the separability problem that appears in quantum computation, can be seen as a normal completion problem where additional constraints need to be met; see [5] for details. In that context, minimizing the size of the matrix corresponds to minimizing the number of states in the separable representation. We will end this paper with considering the problem of completing two matrices to make them commuting.

The paper is organized as follows. In Section 2, we obtain an improved lower bound for $\text{nd}(A)$ by showing that $\text{nd}(A) \geq \max\{i_+([A, A^*]), i_-([A, A^*])\}$, where $i_{\pm}(M)$ denotes the number of positive/negative eigenvalues of the Hermitian matrix M . Using this improved lower bound we are able to show that for some weighted Jordan blocks we have that $\text{nd}(A) = \text{ud}(A)$, providing new evidence for Conjecture 1.1. Next, in Section 3 we examine matrices for which $\text{nd}(A) = 1$. Finally in Section 4 we explore the commuting completion problem.

2 MAIN RESULT AND NORMAL DEFECT CONJECTURE

In this section we shall prove our main result and provide evidence for Conjecture 1.1.

Theorem 2.1 Given $A \in \mathbb{C}^{n \times n}$. Then $\text{nd}(A) \geq \max\{i_+([A, A^*]), i_-([A, A^*])\}$.

Proof. Let $A_{\text{ext}} := \begin{pmatrix} A & A_{12} \\ A_{21} & A_{22} \end{pmatrix}$ be a normal completion of size $(n+q) \times (n+q)$. Denote $a = i_+([A, A^*])$, $b = i_-([A, A^*])$ and $k = n - a - b$, which corresponds to the dimension of the kernel of $[A, A^*]$. Since $[A, A^*]$ is Hermitian, we see that

$$[A, A^*] = U^* \Lambda U, \quad (2.1)$$

for some unitary U and diagonal Λ . Rewriting (2.1) in an alternative form we arrive at

$$U[A, A^*]U^* = [UAU^*, UA^*U^*] = \Lambda. \quad (2.2)$$

Let now

$$B_{\text{ext}} := \begin{pmatrix} B & B_{12} \\ B_{21} & B_{22} \end{pmatrix} = \begin{pmatrix} U & 0 \\ 0 & I \end{pmatrix} \begin{pmatrix} A & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} U^* & 0 \\ 0 & I \end{pmatrix} = \begin{pmatrix} UAU^* & UA_{12} \\ A_{21}U^* & A_{22} \end{pmatrix}.$$

Since A_{ext} is normal, B_{ext} is normal as well. Since B_{ext} is normal, the following equation results from the (1,1) entries of $B_{\text{ext}}B_{\text{ext}}^*$ and $B_{\text{ext}}^*B_{\text{ext}}$ being equal: $BB^* + B_{12}B_{12}^* = B^*B + B_{21}^*B_{21}$. Therefore $\Lambda = [B, B^*] = B_{21}^*B_{21} - B_{12}B_{12}^*$. Writing the above equation in a different form yields

$$\Lambda = [B, B^*] = \begin{pmatrix} B_{21}^* & B_{12} \end{pmatrix} \begin{pmatrix} I_q & 0 \\ 0 & -I_q \end{pmatrix} \begin{pmatrix} B_{21} \\ B_{12}^* \end{pmatrix}. \quad (2.3)$$

Let us remove all rows and columns in Λ that correspond to zeros on the diagonal of Λ giving $\tilde{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_{p-k})$. Next remove the corresponding columns from $\begin{pmatrix} B_{21} \\ B_{12}^* \end{pmatrix}$ and rows from $\begin{pmatrix} B_{21}^* & B_{12} \end{pmatrix}$. This process yields a new equation

$$\tilde{\Lambda} = \begin{pmatrix} \alpha^* & \beta^* \end{pmatrix} \begin{pmatrix} I_q & 0 \\ 0 & -I_q \end{pmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix}, \quad (2.4)$$

where α and β are $q \times (n-k)$.

Since $\tilde{\Lambda}$ is invertible, $\begin{pmatrix} \alpha \\ \beta \end{pmatrix}$ is comprised of linearly independent columns. Thus we can find γ and δ such that $\begin{pmatrix} \alpha & \gamma \\ \beta & \delta \end{pmatrix}$ is invertible. Now we get

$$W := \begin{pmatrix} \tilde{\Lambda} & \star \\ \star & \star \end{pmatrix} = \begin{pmatrix} \alpha^* & \beta^* \\ \gamma^* & \delta^* \end{pmatrix} \begin{pmatrix} I_q & 0 \\ 0 & -I_q \end{pmatrix} \begin{pmatrix} \alpha & \gamma \\ \beta & \delta \end{pmatrix}.$$

This gives that W has q positive eigenvalues and q negative eigenvalues; see e.g. Theorem 4.5.8 in [3]. Let us denote them as follows $\mu_1 \leq \dots \leq \mu_q < 0 < \mu_{q+1} \leq \dots \leq \mu_{2q}$. If we denote the eigenvalues of $\tilde{\Lambda}$ as $\lambda_1, \dots, \lambda_{n-k}$ we see that by the interlacing theorem for Hermitian matrices $\mu_l \leq \lambda_l \leq \mu_{l+2q-n-k}$, where $1 \leq l \leq n-k$; see e.g. Theorem 4.3.15 in [3]. Thus $\tilde{\Lambda}$ has at most q positive eigenvalues and at most q negative eigenvalues. From the way $\tilde{\Lambda}$ was constructed from Λ , we see that Λ has the same number of positive eigenvalues and negative eigenvalues as $\tilde{\Lambda}$ but also has k zero eigenvalues. From equation (2.1) we see that $[A, A^*]$ and Λ have the same inertia (use, e.g., Theorem 4.5.8 in [3]). Thus, $i_+([A, A^*]) = a$, $i_-([A, A^*]) = b$, $k = n - a - b$, $0 \leq a \leq q$ and $0 \leq b \leq q$. This yields $\text{nd}(A) \geq \max\{i_+([A, A^*]), i_-([A, A^*])\}$. \square

Using the well know connection between normal matrices N and pairs of commuting Hermitian matrices $(\text{Re}N, \text{Im}N)$, where $\text{Re}N = \frac{1}{2}(N + N^*)$ and $\text{Im}N = \frac{1}{2i}(N - N^*)$ one can easily deduce the following corollary.

Corollary 2.2 *Let Hermitian matrices $A, B \in \mathbb{C}^{n \times n}$ be given. If there exist Hermitian matrices $A_{ext} = \begin{pmatrix} A & * \\ * & * \end{pmatrix}$, $B_{ext} = \begin{pmatrix} B & * \\ * & * \end{pmatrix}$ of size $(n+q) \times (n+q)$ that commute then $q \geq \max\{i_+(i(BA - AB)), i_-(i(BA - AB))\}$*

Proof. Let $N = \text{Re}N + i\text{Im}N = A + iB$. Calculating NN^* we get $NN^* = (A + iB)(A - iB) = A^2 - iAB + iBA + B^2$. Now if we calculate N^*N we get $N^*N = A^2 + iAB - iBA + B^2$. Thus $NN^* - N^*N = 2i(BA - AB)$. It then follows from Theorem 2.1 that $q \geq \max\{i_+(i(BA - AB)), i_-(i(BA - AB))\}$. \square

Next we explore a class of matrices for which the equality $\text{nd}(A) = \text{ud}(A)$ is true.

Proposition 2.3 *Given $A \in \mathbb{C}^{n \times n}$. Let A be of the form $A := \begin{pmatrix} 0 & a_1 & & 0 \\ & \ddots & \ddots & \\ & & & a_{n-1} \\ 0 & & & 0 \end{pmatrix}$ with either $|a_1| = \dots = |a_l| > \dots > |a_{n-1}| > 0$ or $0 < |a_1| < \dots < |a_{n-l}| = \dots = |a_{n-1}|$, where $1 \leq l \leq n-1$. Then $\text{nd}(A) = \text{ud}(A) = n-l$.*

Proof. Let $\alpha := |a_1| = \dots = |a_l| = \|A\|$, where $\|\cdot\|$ denotes the spectral norm. By Proposition 5.4 in [7] we have that $\text{nd}(A) \leq \text{ud}(A)$. Recall that $\text{ud}(A) = \text{rank}(\|A\|^2 - A^*A)$.

Thus

$$\begin{aligned} \text{ud}(A) &= \text{rank} \left\{ \begin{pmatrix} \alpha^2 & & 0 \\ & \ddots & \\ 0 & & \alpha^2 \end{pmatrix} - \begin{pmatrix} 0 & & 0 \\ & |a_1|^2 & \\ 0 & & \ddots & \\ & & & |a_{n-1}|^2 \end{pmatrix} \right\} \\ &= \text{rank} \begin{pmatrix} \alpha^2 & & & & 0 \\ & 0 & & & \\ & & \ddots & & \\ & & & 0 & \\ & & & & \alpha^2 - |a_{l+1}|^2 \\ & & & & \ddots \\ 0 & & & & & \alpha^2 - |a_{n-1}|^2 \end{pmatrix} \\ &= n - l. \end{aligned}$$

Thus we have that the $\text{nd}(A) \leq \text{ud}(A) = n - l$.

To achieve a lower bound for $\text{nd}(A)$, we notice that Theorem 2.1 yields $\text{nd}(A) \geq \max\{i_+([A, A^*]), i_-([A, A^*])\}$. Observe that

$$\begin{aligned} [A, A^*] &= \begin{pmatrix} |a_1|^2 & & & & & & & & & & 0 \\ & |a_2|^2 - |a_1|^2 & & & & & & & & & \\ & & \ddots & & & & & & & & \\ & & & |a_l|^2 - |a_{l-1}|^2 & & & & & & & \\ & & & & |a_{l+1}|^2 - |a_l|^2 & & & & & & \\ & & & & & \ddots & & & & & \\ & & & & & & |a_{n-1}|^2 - |a_{n-2}|^2 & & & & \\ 0 & & & & & & & & & & -|a_{n-1}|^2 \end{pmatrix} \\ &= \begin{pmatrix} \alpha^2 & & & & & & & & & & 0 \\ & 0 & & & & & & & & & \\ & & \ddots & & & & & & & & \\ & & & 0 & & & & & & & \\ & & & & |a_{l+1}|^2 - \alpha^2 & & & & & & \\ & & & & & \ddots & & & & & \\ 0 & & & & & & |a_{n-1}|^2 - |a_{n-2}|^2 & & & & -|a_{n-1}|^2 \end{pmatrix}. \end{aligned}$$

Therefore $i_+([A, A^*]) = 1$ and $i_-([A, A^*]) = n - l$. Clearly $n - l \geq 1$, thus $\text{nd}(A) \geq n - l$. We now have $\text{nd}(A) \leq n - l$ and $\text{nd}(A) \geq n - l$. We conclude that $\text{nd}(A) = \text{ud}(A) = n - l$. The proof for the second part of the statement is similar. \square

We observe that the matrices in Proposition 2.3 are unitarily irreducible, as the following lemma shows.

Lemma 2.4 *Given $A \in \mathbb{C}^{n \times n}$. Let A be strictly upper triangular with $a_{i,i+1} \neq 0$ for $1 \leq i \leq n - 1$. Then A is unitarily irreducible.*

Proof. Let $A = U^* \begin{pmatrix} A_1 & 0 \\ 0 & A_2 \end{pmatrix} U$ with U unitary, $A_1 \in \mathbb{C}^{p \times p}$ and $A_2 \in \mathbb{C}^{q \times q}$ with $p + q = n$. We need to show that $p = 0$ or $q = 0$. Notice that $n - 1 = \text{rank}(A) = \text{rank}(A_1) + \text{rank}(A_2) \leq p + q = n$. So there exist two possibilities, either $\text{rank}(A_1) = p$ and $\text{rank}(A_2) = q - 1$ or $\text{rank}(A_1) = p - 1$ and $\text{rank}(A_2) = q$. Without loss of generality we assume that $\text{rank}(A_1) = p$ and $\text{rank}(A_2) = q - 1$. Thus A_1 is invertible and $\text{rank}(A_1^k) = p$ for all k . But $0 = \text{rank}(A^n) \geq \text{rank}(A_1^n) = p$. It follows that $p = 0$. \square

3 WHEN THE NORMAL DEFECT EQUALS ONE

Inspired by [4] we consider the case when the normal defect is equal to one. For the class of matrices considered before we have the following observation.

Proposition 3.1 *Given a matrix A of the form $A := \begin{pmatrix} 0 & a_1 & & 0 \\ & \ddots & \ddots & \\ & & & a_{n-1} \\ 0 & & & 0 \end{pmatrix}$, with $a_1, \dots, a_{n-1} \in \mathbb{C} \setminus \{0\}$. Then $\text{nd}(A) = 1$ if and only if $|a_1| = \dots = |a_{n-1}| =: \alpha$. Furthermore, when $n \geq 4$ all minimal normal completions have the form $A_{\text{ext}} := \begin{pmatrix} 0 & a_1 & & 0 & 0 \\ & \ddots & \ddots & & \vdots \\ & & & a_{n-1} & 0 \\ 0 & & & 0 & a_n \\ a_{n+1} & 0 & \dots & & 0 \end{pmatrix}$, with $\alpha = |a_n| = |a_{n+1}|$.*

Proof. Observe that

$$[A, A^*] = \begin{pmatrix} |a_1|^2 & & & & 0 \\ & |a_2|^2 - |a_1|^2 & & & \\ & & \ddots & & \\ & & & |a_{n-1}|^2 - |a_{n-2}|^2 & \\ 0 & & & & -|a_{n-1}|^2 \end{pmatrix},$$

which yields that $\max(i_+([A, A^*]), i_-([A, A^*])) \geq 1$, and equality holds if and only if $|a_1| = \dots = |a_{n-1}|$. Using Theorem 2.1 it follows that for $\text{nd}(A) = 1$ it is necessary that $|a_1| = \dots = |a_{n-1}|$.

Let now $|a_1| = \dots = |a_{n-1}|$, and put

$$A_{\text{ext}} := \begin{pmatrix} 0 & a_1 & & 0 & \beta_1 \\ & \ddots & \ddots & & \beta_2 \\ & & & a_{n-1} & \vdots \\ 0 & & & 0 & \beta_n \\ \gamma_1 & \gamma_2 & \dots & \gamma_n & \delta \end{pmatrix} = \begin{pmatrix} A & \beta \\ \gamma & \delta \end{pmatrix},$$

where $\beta := (\beta_1 \ \dots \ \beta_n)^T$ and $\gamma := (\gamma_1 \ \dots \ \gamma_n)$. Let us assume that A_{ext} is normal.

Then we see that $AA^* + \beta\beta^* = A^*A + \gamma^*\gamma$, or equivalently $[A, A^*] = \gamma^*\gamma - \beta\beta^*$. Recall that

$$\begin{aligned} [A, A^*] &= \begin{pmatrix} |a_1|^2 & & & & 0 \\ & 0 & & & \\ & & \ddots & & \\ & & & 0 & \\ 0 & & & & -|a_{n-1}|^2 \end{pmatrix} \\ &= \begin{pmatrix} \bar{\gamma}_1 \\ \vdots \\ \bar{\gamma}_n \end{pmatrix} (\gamma_1 \ \dots \ \gamma_n) - \begin{pmatrix} \beta_1 \\ \vdots \\ \beta_n \end{pmatrix} (\bar{\beta}_1 \ \dots \ \bar{\beta}_n) \\ &= \begin{pmatrix} |\gamma_1|^2 - |\beta_1|^2 & \bar{\gamma}_1\gamma_2 - \beta_1\bar{\beta}_2 & \dots & \bar{\gamma}_1\gamma_n - \beta_1\bar{\beta}_n \\ \bar{\gamma}_2\gamma_1 - \beta_2\bar{\beta}_1 & |\gamma_2|^2 - |\beta_2|^2 & \dots & \bar{\gamma}_2\gamma_n - \beta_2\bar{\beta}_n \\ \vdots & \vdots & \ddots & \vdots \\ \bar{\gamma}_n\gamma_1 - \beta_n\bar{\beta}_1 & \bar{\gamma}_n\gamma_2 - \beta_n\bar{\beta}_2 & \dots & |\gamma_n|^2 - |\beta_n|^2 \end{pmatrix} =: W = (w_{ij})_{i,j=1}^n. \end{aligned}$$

As $|a_1|, |a_{n-1}| > 0$ we get that $|\gamma_1|^2 = |a_1|^2 + |\beta_1|^2 \neq 0$, $|\beta_n|^2 = |a_{n-1}|^2 + |\gamma_n|^2 \neq 0$. From $w_{i,i} = 0$, for $2 \leq i \leq n-1$, we see that $|\gamma_i| = |\beta_i|$ for $i = 2, \dots, n-1$. From $w_{1,i} = 0$ for $i = 2, \dots, n-1$, we see that $\bar{\gamma}_i\gamma_1 - \beta_i\bar{\beta}_1 = 0$. This implies that $|\gamma_i||\gamma_1| - |\beta_i||\beta_1| = 0$.

Since $|\gamma_i| = |\beta_i|$ for $i = 2, \dots, n-1$, we get $|\gamma_i|(|\gamma_1| - |\beta_1|) = 0$. So either $|\gamma_i| = |\beta_i| = 0$ or $|\gamma_1| = |\beta_1|$. But $|\gamma_1| \neq |\beta_1|$ since $|\gamma_1| = |a_1|^2 + |\beta_1|^2 > |\beta_1|^2$. Therefore $|\gamma_i| = |\beta_i| = 0$ for $i = 2, \dots, n-1$.

To find $\beta_1, \beta_n, \gamma_1$, and γ_n we observe the following equation that results from A_{ext} being normal, $A\gamma^* + \beta\bar{\delta} = A^*\beta + \gamma^*\delta$. Rewriting we see

$$\begin{pmatrix} 0 & a_1 & & 0 \\ & \ddots & \ddots & \\ & & & a_{n-1} \\ 0 & & & 0 \end{pmatrix} \begin{pmatrix} \bar{\gamma}_1 \\ 0 \\ \vdots \\ 0 \\ \bar{\gamma}_n \end{pmatrix} + \begin{pmatrix} \beta_1 \\ 0 \\ \vdots \\ 0 \\ \beta_n \end{pmatrix} \bar{\delta} = \begin{pmatrix} 0 & & & 0 \\ \bar{a}_1 & \ddots & & \\ & \ddots & \ddots & \\ 0 & & & \bar{a}_{n-1} & 0 \end{pmatrix} \begin{pmatrix} \beta_1 \\ 0 \\ \vdots \\ 0 \\ \beta_n \end{pmatrix} + \begin{pmatrix} \bar{\gamma}_1 \\ 0 \\ \vdots \\ 0 \\ \bar{\gamma}_n \end{pmatrix} \delta. \text{ Simplifying we see that } \begin{pmatrix} 0 \\ \vdots \\ 0 \\ a_{n-1}\bar{\gamma}_n \\ 0 \end{pmatrix} + \begin{pmatrix} \beta_1\bar{\delta} \\ 0 \\ \vdots \\ 0 \\ \beta_n\bar{\delta} \end{pmatrix} = \begin{pmatrix} 0 \\ \bar{a}_1\beta_1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} + \begin{pmatrix} \bar{\gamma}_1\delta \\ 0 \\ \vdots \\ 0 \\ \bar{\gamma}_n\delta \end{pmatrix}. \text{ As } n \geq 4, \text{ we see that}$$

$$\beta_1\bar{\delta} = \bar{\gamma}_1\delta, \beta_n\bar{\delta} = \bar{\gamma}_n\delta \quad (3.1)$$

$$\bar{a}_1\beta_1 = 0, \quad (3.2)$$

$$a_{n-1}\bar{\gamma}_n = 0. \quad (3.3)$$

From (3.2) we get $|\beta_1| = 0$ since $|a_1| \neq 0$. From (3.3) we get $|\gamma_n| = 0$. From (3.1) we get $|\delta| = 0$ since $\beta_1 = 0$ and $\gamma_1 \neq 0$. Finally from the equations for $w_{1,1}$ and $w_{n,n}$ we see that $|\gamma_1| = |\beta_n| = \alpha$. In conclusion, for $n \geq 4$ normality of A_{ext} implies that it is of the form

$$A_{ext} := \begin{pmatrix} 0 & a_1 & & 0 & 0 \\ & \ddots & \ddots & & \vdots \\ & & & a_{n-1} & 0 \\ 0 & & & 0 & a_n \\ a_{n+1} & 0 & \dots & & 0 \end{pmatrix}, |a_1| = \dots = |a_{n+1}| \neq 0. \quad (3.4)$$

Clearly, for any n , if A_{ext} is as in (3.4), then A_{ext} is normal. This proves the proposition. \square

It is worth mentioning that for $n = 2$ or $n = 3$ the above proposition does not describe the full situation. For example, consider $A = \begin{pmatrix} 0 & \sqrt{3} \\ 0 & 0 \end{pmatrix}$ and $\tilde{A} := \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$.

Then $A_{ext} := \begin{pmatrix} 0 & \sqrt{3} & 1 \\ 0 & 0 & 2 \\ 2 & 1 & \sqrt{3} \end{pmatrix}$ and $\tilde{A}_{ext} := \begin{pmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & \sqrt{2} \\ \sqrt{2} & 0 & 1 & 0 \end{pmatrix}$ also yield minimal normal completions for A and \tilde{A} , respectively.

Next we determine the eigenvalues of the completed normal matrix A_{ext} from Proposition 3.1.

Lemma 3.2 For the normal matrix $A_{ext} := \begin{pmatrix} 0 & |\alpha|e^{i\theta_1} & & 0 & 0 \\ & \ddots & \ddots & & \vdots \\ & & & |\alpha|e^{i\theta_{n-1}} & 0 \\ 0 & & & 0 & |\alpha|e^{i\theta_n} \\ |\alpha|e^{i\theta_{n+1}} & 0 & \dots & & 0 \end{pmatrix}$ all eigenvalues of A_{ext} are exactly $\lambda_k = |\alpha|e^{i\frac{(\psi+2\pi k)}{n+1}}$ where $\psi = \theta_1 + \dots + \theta_{n+1}$ and $k = 0, \dots, n$.

Proof. Computing the characteristic polynomial of A_{ext} we see that $\lambda^{n+1} = |\alpha|^{n+1}e^{i\psi}$, where $\psi = \theta_1 + \dots + \theta_{n+1}$. Thus $\lambda_k = |\alpha|e^{i\frac{(\psi+2\pi k)}{n+1}}$ with $k = 0, \dots, n$. \square

In [4] the following problem was considered. Let $\lambda_0, \dots, \lambda_n$ and μ_1, \dots, μ_n be two sequences of complex numbers. When can one find a $(n+1) \times (n+1)$ normal matrix with eigenvalues $\lambda_0, \dots, \lambda_n$ whose $n \times n$ principal submatrix has eigenvalues μ_1, \dots, μ_n ? In [4] the author obtained the following result. Define $\Delta(\lambda) := \frac{\prod_{j=1}^n (\lambda - \mu_j)}{\prod_{k=0}^n (\lambda - \lambda_k)}$. Then there exists a normal matrix A with spectrum $\lambda_0, \dots, \lambda_n$ and with a $n \times n$ principal submatrix with spectrum μ_1, \dots, μ_n , if and only if the rational function Δ has only simple poles and $\text{Res}_{\lambda_k}(\Delta(\lambda)) \geq 0$, $k \in 0, \dots, n$. If we take A_{ext} in Lemma 3.2 we observe that A_{ext} has spectrum $\lambda_k = |\alpha|e^{i\frac{\psi+2\pi k}{n+1}}$, where $\psi = \theta_1 + \dots + \theta_{n+1}$ and $k = 0, \dots, n$. When we remove the last row and last column in A_{ext} the remaining matrix has eigenvalues $\mu_i = 0$, $i = 1, \dots, n$. Thus by the result in [4] we must have that $\text{Res}_{\lambda_k}(\frac{\lambda^n}{\prod_{k=0}^n (\lambda - \lambda_k)}) \geq 0$ is satisfied. This is indeed true, since $\text{Res}_{\lambda_k}(\frac{\lambda^n}{\lambda^{n+1} - |\alpha|^{n+1}e^{i\psi}}) = \lim_{\lambda \rightarrow \lambda_k} \frac{(\lambda - \lambda_k)\lambda^n}{\lambda^{n+1} - |\alpha|^{n+1}e^{i\psi}} = \lim_{\lambda \rightarrow \lambda_k} \frac{\lambda^n + (\lambda - \lambda_k)n\lambda^{n-1}}{(n+1)\lambda^n} = \frac{1}{n+1} \geq 0$, where we used L'Hôpital's rule in the second equality.

4 COMMUTING DEFECT

In [1] the following Minimal Commuting Completion problem was introduced. Given $A_1, \dots, A_d \in \mathbb{C}^{n \times n}$, how do we find matrices $(A_1)_{ext}, \dots, (A_d)_{ext} \in \mathbb{C}^{(n+q) \times (n+q)}$ of smallest possible size with

$$(A_i)_{ext} = \begin{pmatrix} A_i & * \\ * & * \end{pmatrix}$$

such that $[(A_i)_{ext}, (A_j)_{ext}] = 0$, $i \neq j$. We will restrict our investigation to completing only two matrices: given $A, B \in \mathbb{C}^{n \times n}$ how do we find A_{ext}, B_{ext} of smallest possible size, with

$$A_{ext} = \begin{pmatrix} A & A_{12} \\ A_{21} & A_{22} \end{pmatrix}, B_{ext} = \begin{pmatrix} B & B_{12} \\ B_{21} & B_{22} \end{pmatrix}, \quad (4.1)$$

such that $[A_{ext}, B_{ext}] = 0$. We shall call the smallest possible number q the *commuting defect*, and denote it $\text{cd}(A, B)$. Clearly, $\text{cd}(A, B) = 0$ if and only if $[A, B] = 0$. As shown in [1], $\text{cd}(A, B) \geq \frac{1}{2} \text{rank}([A, B])$. Indeed if $[A_{ext}, B_{ext}] = 0$ then $AB - BA = B_{21}A_{21} - A_{12}B_{12}$ and thus $\text{rank}([A, B]) \leq \text{rank}(B_{21}A_{21}) + \text{rank}(A_{12}B_{12}) \leq q + q = 2q$. The results in [1] also show that $\text{cd}(A, B) \leq n$. One easily sees this by taking $A_{ext} = \begin{pmatrix} A & B \\ B & A \end{pmatrix}$ and $B_{ext} = \begin{pmatrix} B & A \\ A & B \end{pmatrix}$.

One useful observation for this problem is that if two square matrices C and D satisfy $CD = \alpha I$ for some $\alpha \neq 0$, then automatically $CD = DC$. With this in mind we introduce the Minimal Inverse Completion problem: Given $A, B \in \mathbb{C}^{n \times n}$, how do we find $A_{ext}, B_{ext} \in \mathbb{C}^{(n+q) \times (n+q)}$ as in (4.1) of smallest possible size such that $A_{ext}B_{ext} = \alpha I_{n+q}$, for some $\alpha \neq 0$. We shall call this smallest number q the *inverse defect* and denote it by $\text{id}(A, B)$. The inverse defect of a pair of matrices is easily determined as the following theorem shows.

Theorem 4.1 *Given $A, B \in \mathbb{C}^{n \times n}$. Suppose α is the nonzero eigenvalue of AB with the highest geometric multiplicity. Then $\text{id}(A, B) = \text{rank}(\alpha I_n - AB)$.*

Proof. Let $A_{ext} := \begin{pmatrix} A & ? \\ ? & ? \end{pmatrix}$ and $B_{ext} := \begin{pmatrix} B & ? \\ ? & ? \end{pmatrix}$ exist such that $A_{ext}B_{ext} = \alpha I_{n+q}$, where $\alpha \neq 0$. We notice that $A_{ext}B_{ext} = \alpha I_{n+q}$ if and only if $\text{rank} \begin{pmatrix} \alpha I_{n+q} & A_{ext} \\ B_{ext} & I_{n+q} \end{pmatrix} = n + q$. We define $W := \begin{pmatrix} \alpha I_n & 0 & A & ? \\ 0 & \alpha I_q & ? & ? \\ B & ? & I_n & 0 \\ ? & ? & 0 & I_q \end{pmatrix}$. After a permutation similarity we arrive at

$\tilde{W} := \begin{pmatrix} \alpha I_q & 0 & ? & ? \\ 0 & \alpha I_n & A & ? \\ ? & B & I_n & 0 \\ ? & ? & 0 & I_q \end{pmatrix}$. Therefore we would to find a minimal rank completion \tilde{W}_{compl}

of \tilde{W} so that $\text{rank}(\tilde{W}_{compl}) = n + q$. Fixing q , we must minimize $\text{rank}(\tilde{W})$. By inspection we see that \tilde{W} is a partial banded matrix with a pattern J . Therefore we can apply Theorem 1.1 from [6]. We now have that $\text{minrank}(\tilde{W}) = \max_{T \subset J}(\text{minrank}W_T)$, where W_T is the partial matrix obtained from \tilde{W} by only keeping the known entries that lie in the triangular subpattern T . From this it follows that $\text{minrank}(\tilde{W}) = \max\{n + q, \text{rank}\begin{pmatrix} \alpha I_n & A \\ B & I_n \end{pmatrix}\}$. Notice that $\text{rank}\begin{pmatrix} \alpha I_n & A \\ B & I_n \end{pmatrix} = n + \text{rank}(\alpha I_n - AB)$ by a Schur complement argument. Now suppose $q < \text{rank}(\alpha I_n - AB)$. Then $\text{minrank}(\tilde{W}) = \max\{n + q, \text{rank}\begin{pmatrix} \alpha I_n & A \\ B & I_n \end{pmatrix}\} = n + \text{rank}(\alpha I_n - AB) > n + q$. So then \tilde{W} does not have a completion of size $n + q$, so we cannot find A_{ext}, B_{ext} of size $(n + q) \times (n + q)$ so that $A_{ext}B_{ext} = \alpha I_{n+q}$. Thus $\text{id}(A, B) \geq \text{rank}(\alpha I_n - AB)$.

Now let $\alpha \neq 0$ be a nonzero eigenvalue of AB with highest geometric multiplicity (or equivalently, the $\alpha \neq 0$ so that $\text{rank}(\alpha I_n - AB)$ is minimal), and set $q = \text{rank}(\alpha I_n - AB)$. Then $\text{minrank}\tilde{W} = \max\{n + q, \text{rank}\begin{pmatrix} \alpha I_n & A \\ B & I_n \end{pmatrix}\} = n + q$. Thus \tilde{W} has a completion of rank $n + q$ and consequently we can find A_{ext} and B_{ext} such that $A_{ext}B_{ext} = \alpha I_{n+q}$. Thus $\text{id}(A, B) \leq \text{rank}(\alpha I_n - AB)$. Therefore $\text{id}(A, B) \leq \text{rank}(\alpha I_n - AB)$ together with $\text{id}(A, B) \geq \text{rank}(\alpha I_n - AB)$ yields $\text{id}(A, B) = \text{rank}(\alpha I_n - AB)$. \square

Let us outline how to find a minimal inverse completion.

Algorithm 4.2 Let A, B and α be as in Theorem 4.1 and put $q = \text{id}(A, B)$ and $p = n - q$. First determine an invertible matrix S so that $SABS^{-1} = \begin{pmatrix} \alpha I_p & P \\ 0 & Q \end{pmatrix}$ for some P and Q of sizes $p \times q$ and $q \times q$, respectively. Write now $SA = \begin{pmatrix} C_{11} & C_{22} \\ C_{21} & C_{22} \end{pmatrix}$, $BS^{-1} = \begin{pmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{pmatrix}$ with C_{11}, D_{11} of size $p \times p$ and C_{22}, D_{22} of size $q \times q$. Choose X and Y to be invertible matrices of size $q \times q$ and let

$$D_{ext} = \begin{pmatrix} D_{11} & D_{12} & -(D_{11}C_{12} + D_{12}C_{22})Y^{-1} \\ D_{21} & D_{22} & -(D_{21}C_{12} + D_{22}C_{22} - \alpha I)Y^{-1} \\ 0 & X & -XC_{22}Y^{-1} \end{pmatrix}$$

and put $B_{ext} = D_{ext}(S \oplus I_q)$, $A_{ext} = \alpha B_{ext}^{-1}$. Then one can check that A_{ext} and B_{ext} have the required form and (obviously) $A_{ext}B_{ext} = \alpha I$.

Let us try the algorithm out on an example.

Example 4.3 Let $A = \begin{pmatrix} 1 & 0 \\ 0 & 2 \end{pmatrix}$ and $B = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$. Then $AB = \begin{pmatrix} 0 & 1 \\ 2 & 0 \end{pmatrix}$ has eigenvalues $\pm\sqrt{2}$. Let us choose $\alpha = \sqrt{2}$ and $S = \begin{pmatrix} \frac{1}{2} & \frac{1}{2\sqrt{2}} \\ \frac{1}{2} & -\frac{1}{2\sqrt{2}} \end{pmatrix}$. Then $BS^{-1} = \begin{pmatrix} \sqrt{2} & -\sqrt{2} \\ 1 & 1 \end{pmatrix}$, $SA = \begin{pmatrix} \frac{1}{2} & \frac{1}{\sqrt{2}} \\ \frac{1}{2} & -\frac{1}{\sqrt{2}} \end{pmatrix}$. Thus we get $D_{ext} = \begin{pmatrix} \sqrt{2} & -\sqrt{2} & -\frac{2}{y} \\ 1 & 1 & \frac{\sqrt{2}}{y} \\ 0 & x & \frac{x}{\sqrt{2}y} \end{pmatrix}$ and $A_{ext} = \begin{pmatrix} 1 & 0 & \frac{2\sqrt{2}}{x} \\ 0 & 2 & -\frac{4}{x} \\ -\frac{y}{\sqrt{2}} & y & -\frac{2y}{x} \end{pmatrix}$, $B_{ext} = \begin{pmatrix} 0 & 1 & -\frac{2}{y} \\ 1 & 0 & \frac{\sqrt{2}}{y} \\ \frac{x}{2} & -\frac{x}{2\sqrt{2}} & \frac{x}{\sqrt{2}y} \end{pmatrix}$.

As observed before, we have that $\text{cd}(A, B) \leq \text{id}(A, B)$. In general there is no equality. For instance, when $A = (0)$ and $B = (1)$ we have $\text{cd}(A, B) = 0$ and $\text{id}(A, B) = 1$. That such a simple example exists seems to be due to the fact that we are excluding the possibility $\alpha = 0$ in the definition of $\text{id}(A, B)$; we do this since $CD = 0$ does not imply $DC = 0$. But now one can ask what happens when A and B are nonsingular. Even in that case one can in general improve upon the estimate $\text{cd}(A, B) \leq \text{id}(A, B)$ by doing the following. Suppose there is an invertible matrix S so that

$$SAS^{-1} = \begin{pmatrix} A_1 & & 0 \\ & \ddots & \\ 0 & & A_d \end{pmatrix}, SBS^{-1} = \begin{pmatrix} B_1 & & 0 \\ & \ddots & \\ 0 & & B_d \end{pmatrix} \quad (4.2)$$

where A_i and B_i , $i = 1, \dots, d$, are square matrices of the same nontrivial size. Then completing A_i and B_i to $(A_i)_{ext}$ and $(B_i)_{ext}$ that commute for all $i \in \{1, \dots, d\}$, yields completions A_{ext} and B_{ext} for A and B respectively, that commute. Thus $\text{cd}(A, B) \leq \sum_{i=1}^d \text{cd}(A_i, B_i) \leq \sum_{i=1}^d \text{id}(A_i, B_i)$. We are now led to the following question: Let A and B be nonsingular matrices so that for no invertible S we have that (4.2) holds with $d \geq 2$. Is it then true that $\text{cd}(A, B) = \text{id}(A, B)$?

The questions in this section may also be pursued in the class of real symmetric matrices. In other words, let A and B be real symmetric and look for A_{ext} and B_{ext} that are also real symmetric. As a complex symmetric matrix N is normal if and only if the real symmetric matrices $A = \operatorname{Re} N$ and $B = \operatorname{Im} N$ commute, Corollary 2.2 applies. The real symmetric case is of interest in deriving multivariable quadrature formulas; see [1]. In their setting A and B have a tridiagonal block form and A_{ext} and B_{ext} are required to have this form as well. For this reason it may not be optimal to look for A_{ext} and B_{ext} with $A_{ext}B_{ext} = \alpha I_{n+q}$. We hope to further pursue the symmetric case in a future publication.

ACKNOWLEDGEMENT

Both authors were partially supported by NSF grant DMS-0500678. David P. Kimsey performed the research as part of an REU project.

References

- [1] Ilan Degani, Jeremy Schiff, and David J. Tannor. Commuting extensions and cubature formulae. *Numer. Math.*, 101(3):479–500, 2005.
- [2] P. R. Halmos. Subnormal suboperators and the subdiscrete topology. In *Anniversary volume on approximation theory and functional analysis (Oberwolfach, 1983)*, volume 65 of *Internat. Schriftenreihe Numer. Math.*, pages 49–65. Birkhäuser, Basel, 1984.
- [3] Roger A. Horn and Charles R. Johnson. *Matrix analysis*. Cambridge University Press, Cambridge, 1990. Corrected reprint of the 1985 original.
- [4] S. M. Malamud. Inverse spectral problem for normal matrices and the Gauss-Lucas theorem. *Trans. Amer. Math. Soc.*, 357(10):4043–4064 (electronic), 2005.
- [5] H. J. Woerdeman. The separability problem and normal completions. *Linear Algebra Appl.*, 376:85–95, 2004.
- [6] Hugo J. Woerdeman. Minimal rank completions of partial banded matrices. *Linear and Multilinear Algebra*, 36(1):59–68, 1993.
- [7] Hugo J. Woerdeman. Hermitian and normal completions. *Linear and Multilinear Algebra*, 42(3):239–280, 1997.